# CUDD <br> Colorado University Decision Diagram Package 

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(3) Basic Functions

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## CUDD

- CUDD is the Colorado University Decision Diagram package.
- It is a $\mathrm{C} / \mathrm{C}++$ library for creating the following different types of decision diagrams:
- BDD: Binary Decision Diagrams.
- ZDD: Zero-Suppressed BDDs.
- ADD: Algebraic Decision Diagrams.
- The slides, source code, and all documentation related to this lecture are available here:
https://github.com/luigicapogrosso/SDL


## Getting CUDD

- The CUDD package is available via anonymous FTP from vlsi.colorado.edu.
- You can download the CUDD package from the server using an FTP client such as FileZilla or you can use the ftp command from the command line.
- Alternatively, you can download the latest version of CUDD directly from the SDL GitHub repository, so:
\$ git clone
https://github.com/luigicapogrosso/SDL.git


## Getting CUDD (cont'd)

- The library is tested using GCC (9.4.0) and GNU Make (9.4.0). To build the library from sources in a clean way, it is preferable that you set up a build subdirectory, say:
\$ cd SDL/lecture_01/cudd-3.0.0
\$ mkdir objdir \&\& cd objdir
\$ ../configure --prefix=\$HOME/<path>
\$ make \&\& make install


## Including and linking the CUDD library

- To build an application that uses the CUDD package, you should add, in your source code, the following lines:
- \#include "cudd.h"
- \#include "util.h"
- To compile and link a C program that uses CUDD:
\$ gcc -o main main.c -lcudd -lutil
- Or, you can refer to the following Makefile:
https://github.com/luigicapogrosso/SDL/blob/
master/lecture_01/code/Makefile


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## Garbage Collection

- CUDD has a built-in garbage collection system.
- When a BDD is not used anymore, its memory can be reclaimed.
- To facilitate the garbage collector, we need to "reference" and "dereference" each node in our BDD:
- Cudd_Ref (DdNode*) to reference a node.
- Cudd_RecursiveDeref(DdNode*) to dereference a node and all its descendants.


## Complemented arcs

- Each node of a BDD can be:
- A variable with two children.
- A leaf with a constant value.
- The two children of a node are referred to as the "then" child and the "else" child.
- To assign a value to a BDD, we follow "then" and "else" children until we reach a leaf:
- The value of our assignment is the value of the leaf we reach.
- However: "else" children can be complemented:
- When an "else" child is complemented, then we take the complement of the value of the leaf:
$\star$ i.e., if the value of the leaf is 1 and we have traversed an odd number of complement arcs, the value of our assignment is 0 .


## Complemented arcs: example

- out $=x_{0} \bar{x}_{1}$
- "then" arcs are solid.
- Normal "else" arcs are dashed.
- Complemented "else" arcs are dotted.
- The out arc is complemented:

$$
\begin{aligned}
\overline{\text { out }} & =\bar{x}_{0}+x_{1} \\
& =\bar{x}_{0}+x_{0} x_{1}
\end{aligned}
$$



## The DdManager

- The DdManager is the key data structure of CUDD:
- It must be created before calling any other CUDD function.
- It needs to be passed to almost every CUDD function.
- To initialize the DdManager, we use the following function:

```
DdManager * Cudd_Init(
    unsigned int numVars,
    unsigned int numVarsZ,
    unsigned int numSlots,
    unsigned int cacheSize,
    unsigned long maxMemory
        // initial number of BDD variables (i.e., subtables)
);
```


## The DdManager: C code

```
#include<stdio.h>
#include"cudd.h"
int main()
{
    DdManager* manager = Cudd_Init(0, 0,
        CUDD_UNIQUE_SLOTS, CUDD_CACHE_SLOTS, 0);
    if(manager == NULL)
    {
        printf("Error when initalizing CUDD.\n");
        return 1;
    }
    •••
    return 0;
}
```


## The DdNode

- The DdNode is the core building block of BDDs:

```
struct DdNode {
    DdHalfWord index; // Index of the variable reprented by this node
    DdHalfWord ref;
    DdNode * next;
    union {
    CUDD_VALUE_TYPE value; // for constant nodes
    DdChildren kids; // for internal nodes
    } type;
};
```

- index is a unique index for the variable represented by this node.
- It is permanent: if we reorder variables, the idx remains the same.
- ref stores the reference count for this node.
- It is incremented by Cudd_Ref() and decremented by Cudd_Recursive_Deref().


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## BDD of Boolean functions

- Common manipulations of BDDs can be accomplished by calling operators on variables.
- The CUDD package includes Boolean functions that can be used for BDD operations such as: NOT, AND, NAND, OR, NOR, Exclusive-OR, XNOR, and etc.


## BDD for the NOT Boolean function

- For the NOT Boolean function, we use Cudd_Not ().
- The truth table for a NOT:

| $\mathbf{x}_{1}$ | $\mathbf{f}$ |
| :---: | :---: |
| 0 | 1 |
| 1 | 0 |

- Exercise: write the code to build the BDD for the function $f=\neg X_{1}$.


## BDD for the AND Boolean function

- For the AND Boolean function, we use Cudd_bddAnd ().
- The truth table for an AND:

| $\mathbf{x}_{\mathbf{1}}$ | $\mathbf{x}_{\mathbf{2}}$ | $\mathbf{f}$ |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 0 |
| 1 | 0 | 0 |
| 1 | 1 | 1 |

- Exercise: write the code to build the BDD for the function $f=x_{1} \wedge x_{2}$.


## BDD for the NAND Boolean function

- For the NAND Boolean function, we use Cudd_bddNand ().
- The truth table for a NAND:

| $\mathbf{x}_{\mathbf{1}}$ | $\mathbf{x}_{\mathbf{2}}$ | $\mathbf{f}$ |
| :---: | :---: | :---: |
| 0 | 0 | 1 |
| 0 | 1 | 1 |
| 1 | 0 | 1 |
| 1 | 1 | 0 |

- Exercise: write the code to build the BDD for the function $f=\neg\left(x_{1} \wedge x_{2}\right)$.


## BDD for the OR Boolean function

- For the OR Boolean function, we use Cudd_bddOr ().
- The truth table for a logic OR:

| $\mathbf{x}_{\mathbf{1}}$ | $\mathbf{x}_{\mathbf{2}}$ | $\mathbf{f}$ |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 1 |
| 1 | 0 | 1 |
| 1 | 1 | 1 |

- Exercise: write the code to build the BDD for the function $f=x_{1} \vee x_{2}$.


## BDD for the NOR Boolean function

- For the NOR Boolean function, we use Cudd_bddNor () .
- The truth table for a NOR:

| $\mathbf{x}_{1}$ | $\mathbf{x}_{\mathbf{2}}$ | $\mathbf{f}$ |
| :---: | :---: | :---: |
| 0 | 0 | 1 |
| 0 | 1 | 0 |
| 1 | 0 | 0 |
| 1 | 1 | 0 |

- Exercise: write the code to build the BDD for the function $f=\neg\left(x_{1} \vee x_{2}\right)$.


## BDD for Exclusive-OR Boolean function

- For the Exclusive-OR Boolean function, we use Cudd_bddXor().
- The truth table for an Exclusive-OR:

| $\mathbf{x}_{1}$ | $\mathbf{x}_{\mathbf{2}}$ | $\mathbf{f}$ |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 1 |
| 1 | 0 | 1 |
| 1 | 1 | 0 |

- Exercise: write the code to build the BDD for the function $f=x 1 \oplus x 2$.


## BDD for the XNOR Boolean function

- For the XNOR Boolean function, we use Cudd_bddXnor ().
- The truth table for an XNOR:

| $\mathbf{x}_{1}$ | $\mathbf{x}_{\mathbf{2}}$ | $\mathbf{f}$ |
| :---: | :---: | :---: |
| 0 | 0 | 1 |
| 0 | 1 | 0 |
| 1 | 0 | 0 |
| 1 | 1 | 1 |

- Exercise: write the code to build the BDD for the function $f=\neg\left(x_{1} \oplus x_{2}\right)$.


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## The Half-Adder circuit



This is the schematic of a halfadder circuit that we want to compile into an OBDD. It has the following truth table:

| $\mathbf{x}_{\mathbf{1}}$ | $\mathbf{x}_{\mathbf{2}}$ | sum | carry |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 |
| 0 | 1 | 1 | 0 |
| 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | 1 |

## Create the BDD for sum

```
DdNode *x1 = Cudd_bddIthVar(manager, 0);
DdNode *x2 = Cudd__bddIthVar(manager, 1);
DdNode *and1;
and1 = Cudd_bddAnd(manager, x1, Cudd_Not(x2));
Cudd_Ref(and1);
DdNode *and2;
and2 = Cudd__bddAnd(manager, Cudd_Not(x1), x2);
Cudd_Ref(and2);
DdNode *sum;
sum = Cudd_boddOr(manager, and1, and2);
Cudd_Ref(sum);
Cudd_RecursiveDeref(manager, and1);
Cudd_RecursiveDeref(manager, and2);
```

- Exercise: write the code for carry.


## Restricting the BDD

- Restricting a BDD means assigning a truth value to some of the variables.
- The Cudd_bddRestrict () function returns the restricted BDD.

```
DdNode * Cudd_bddRestrict(
    DdManager * manager, // DD manager
    DdNode * BDD, // The BDD to restrict
    DdNode * restrictBy) // The BDD to restrict by.
```

- $B D D$ is the original BDD to restrict.
- restrictBy is the truth assignment of the variables.


## Print the truth table

```
DdNode *restrictBy;
restrictBy = Cudd_bddAnd(manager, x1, Cudd_Not(x2));
Cudd_Ref(restrictBy);
DdNode *testSum;
testSum = Cudd_bddRestrict(manager, sum, restrictBy);
Cudd_Ref(testSum);
DdNode *testCarry;
testCarry = Cudd_bddRestrict(manager, carry, restrictBy);
Cudd_Ref(testCarry);
printf("x1 = 1, x2 = 0: sum = %d, carry = %d\n",
    1 - Cudd_IsComplement (testSum),
    1 - Cudd_IsComplement(testCarry));
Cudd_RecursiveDeref(manager, restrictBy);
Cudd_RecursiveDeref(manager, testSum);
Cudd_RecursiveDeref(manager, testCarry);
```

- Exercise: Write the code for the complete truth table.


## Print the BDD with graphviz

- The function Cudd_DumpDot () dumps the BDD to a file in GraphViz format.
- The . dot file can be converted to a PDF by the command dot:
\$ dot -O -Tpdf half_adder.dot


## Print the BDD: C code

```
char *inputNames[2];
inputNames[0] = "x1";
inputNames[1] = "x2";
char *outputNames[2];
outputNames[0] = "sum";
outputNames[1] = "carry";
DdNode *outputs[2];
outputs[0] = sum;
Cudd_Ref(outputs[0]);
outputs[1] = carry;
Cudd_Ref(outputs[1]);
FILE *f = fopen("half_adder.dot", "w");
Cudd_DumpDot(manager, 2, outputs, inputNames, outputNames, f);
Cudd_RecursiveDeref(manager, outputs[0]);
Cudd_RecursiveDeref(manager, outputs[1]);
fclose(f);
```


## Variable reordering

- The order of variables can have a tremendous effect on the size of BDDs.
- CUDD provides a rich set of tools for reordering BDDs:
- Automatic reordering (using heuristics) when the number of nodes in the BDD passes a certain threshold.
- Manual reordering using different heuristics.
- Manual reordering with a user-specified variable order.
- The function Cudd_ShuffleHeap () is used to define the variable order:

```
int Cudd_ShuffleHeap(
    DdManager * manager, // DD manager
    int * permutation // required variable permutation
)
```


## Exercise: play with the variable order!

- Create the BDD for the function $x_{1} x_{2}+x_{3} x_{4}+x_{5} x_{6}$.
- Try the following variable orders and compare the results:
- $x_{1}<x_{2}<x_{3}<x_{4}<x_{5}<x_{6}$
- $x_{1}<x_{3}<x_{5}<x_{2}<x_{4}<x_{6}$


## HINTS

- int Cudd_ReadPerm(manager, x2->index) returns the position of variable $x 2$ in the order.
- int Cudd_ReadNodeCount (manager) returns the number of nodes in the BDD.


## Converting BDDs to ZDDs

- Many applications first build a set of BDDs and then derive ZDDs from the BDDs.
- These applications should create the manager with 0 ZDD variables and create the BDDs.
- Then they should call Cudd_zddVarsFromBddVars () to create the necessary ZDD variables-whose number is known once the BDDs are available.


## Converting BDDs to ZDDs (cont'd)

- The simplest conversion from BDDs to ZDDs is a simple change of representation, which preserves the functions.
- Simply put, given a BDD for $f$, a ZDD for $f$ is requested. In this case the correspondence between the BDD variables and ZDD variables is one-to-one.
- Hence, Cudd_zddVarsFromBddVars () should be called with the multiplicity parameter equal to 1 .
- The conversion can then be performed by calling: Cudd_zddPortFromBdd().
- The inverse transformation is performed by calling: Cudd_zddPortToBdd().


## The N-Queens problem

- The $\mathbf{N}$-Queens problem is the problem of placing $N$ non-attacking queens on an $N \times N$ chessboard.
- Our implementation of these benchmarks is based on the description of Kunkle10. We construct a ZDD row-by-row to represent whether the row is in a legal state.
- On the accumulated ZDD we then count the number of satisfying assignments.

