

# Introduction

## The Structure of a Compiler

Linguaggi e Compilatori  
Modulo Compilatori

Alessandra Di Pierro  
`alessandra.dipierro@univr.it`



# What is a compiler?

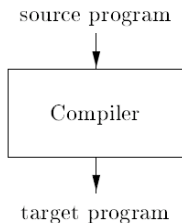


Figure 1.1: A compiler

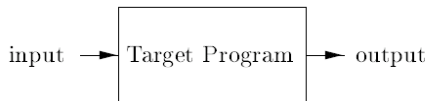


Figure 1.2: Running the target program

## Interpreter

Another kind of language processing

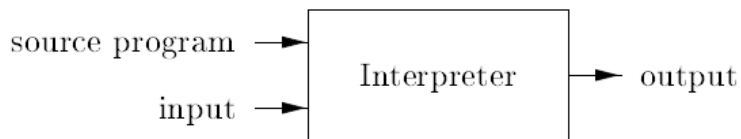


Figure 1.3: An interpreter

## Hybrid Approaches

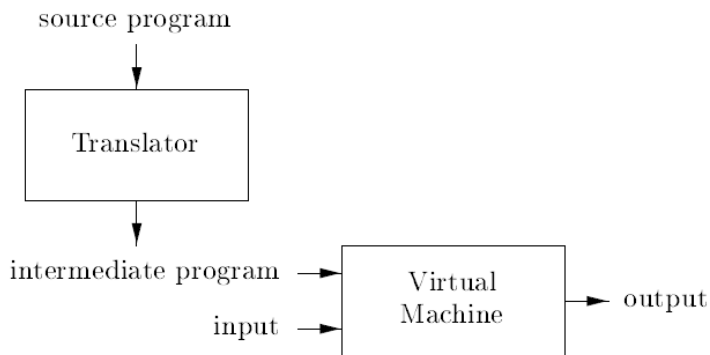


Figure 1.4: A hybrid compiler

# Producing a machine code

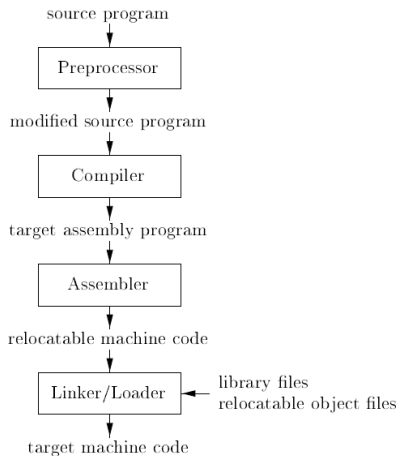


Figure 1.5: A language-processing system

# Phases of a Compiler

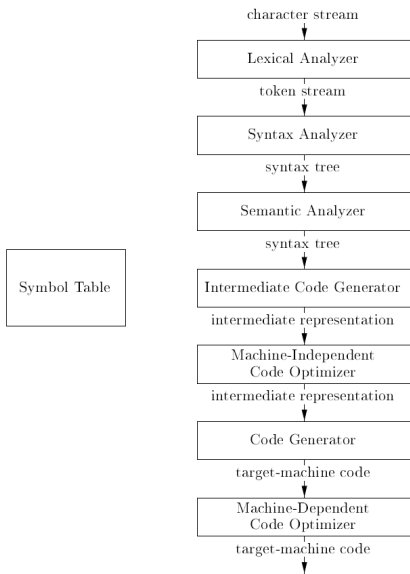
- **Analysis** or *front-end*
- **Synthesis** or *back-end*

The **symbol table** stores information about the entire source program.

Maps variables into attributes, i.e. type, name, dimension, address, etc.

This information helps us detecting inconsistencies and misuses during type checking.

# Compilation process





# Compilation process

1	position	...
2	initial	...
3	rate	...

SYMBOL TABLE

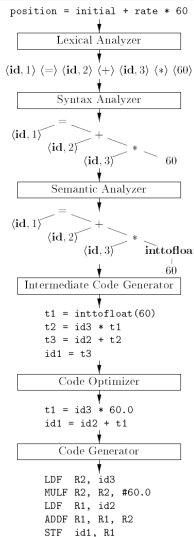


Figure 1.7: Translation of an assignment statement

# Languages and Compilers

The design of programming languages is strongly related to the design of compilers.

Adding new language features places new demand to compilers writers

- **1940's** programs are sequences of 0's and 1's (first electronic computers)
- **Early 1950's** Assembly
- **Late 1950's** Fortran, Cobol, Lisp
- **1970's** C Language
- **1990's** C++, Java

Since 1940's computer architectures has evolved as well!

# High-level Programming Languages

They define programming **abstractions**.

Compilers must translate programs to the target language.

Easier to write programs but the generated target programs run more slowly.

Need for **optimisation**

**Example:** the `register` keyword in C.

## A Simple Example

Objective: to translate programs such as the following simple one:

```
{
    int i; int j; float[100] a; float v; float x;
    while ( true ) {
        do i = i+1; while ( a[i] < v );
        do j = j-1; while ( a[j] > v );
        if ( i >= j ) break;
        x = a[i]; a[i] = a[j]; a[j] = x;
    }
}
```

Figure 2.1: A code fragment to be translated

## A Simple Example (ctd.)

The compiler front end ( [▶ source code](#) ) translates the program into the form:

```
1:  i = i + 1
2:  t1 = a [ i ]
3:  if t1 < v goto 1
4:  j = j - 1
5:  t2 = a [ j ]
6:  if t2 > v goto 4
7:  ifFalse i >= j goto 9
8:  goto 14
9:  x = a [ i ]
10: t3 = a [ j ]
11: a [ i ] = t3
12: a [ j ] = x
13: goto 1
14:
```

Figure 2.2: Simplified intermediate code for the program fragment in Fig. 2.1